

# Rectifiable sets and the Traveling Salesman Problem

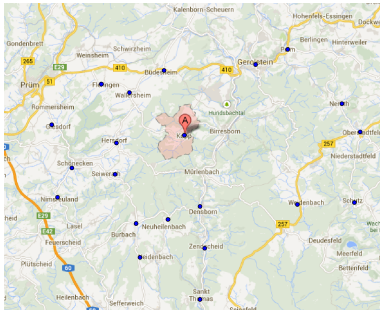
Martí Prats

Universitat Autònoma de Barcelona

July 22, 2013

# First session

# The Traveling Salesman Problem

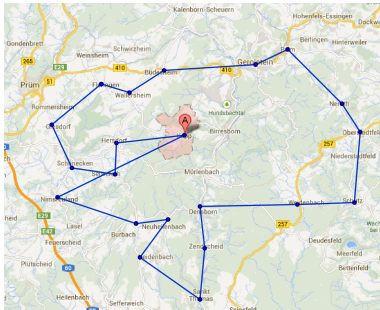


A salesman wants to visit a number of villages and then go back home.



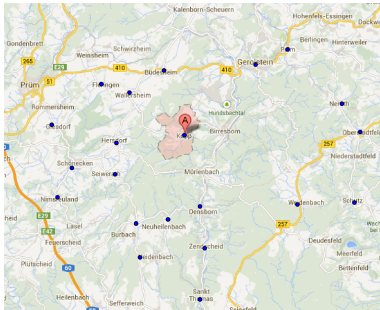


# The Traveling Salesman Problem



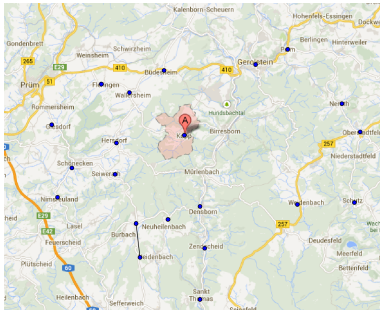
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He wants to find the shortest cycle!

# The Traveling Salesman Problem



The greedy algorithm gives us the minimal spanning tree: choose a vertex.

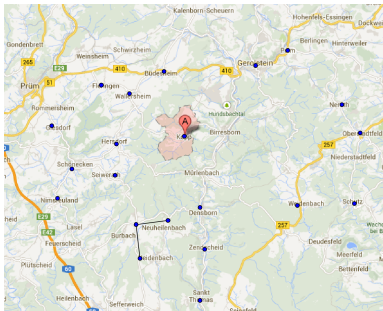
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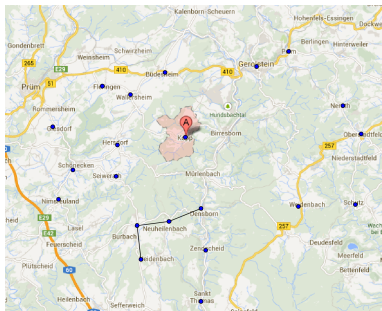


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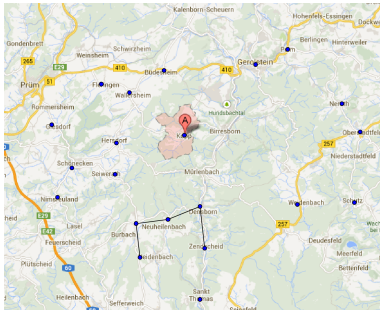


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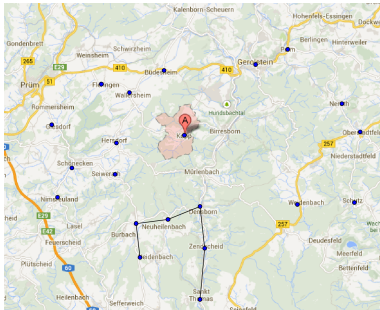


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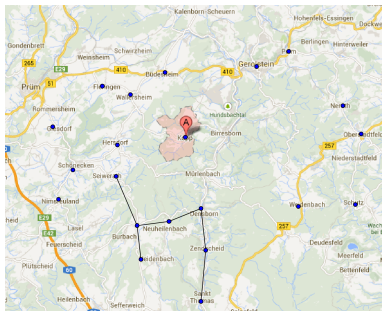


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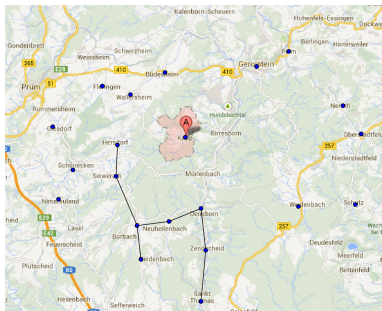
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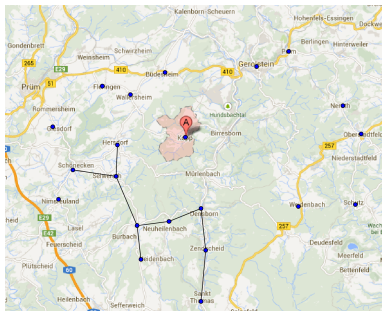


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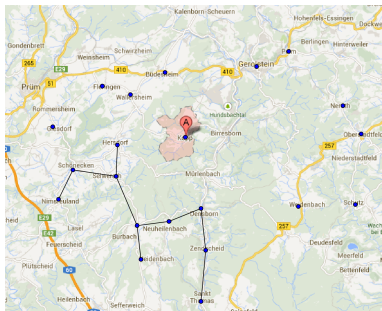


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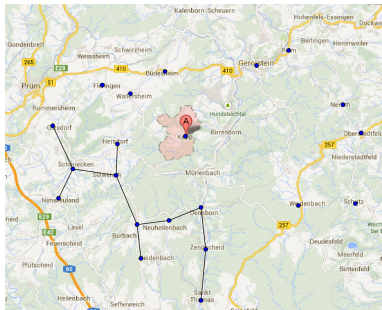


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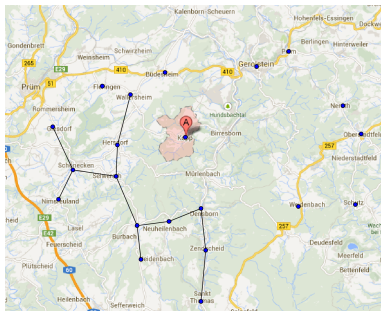


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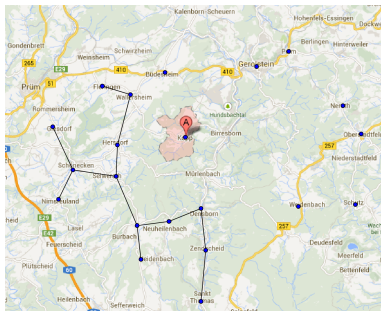


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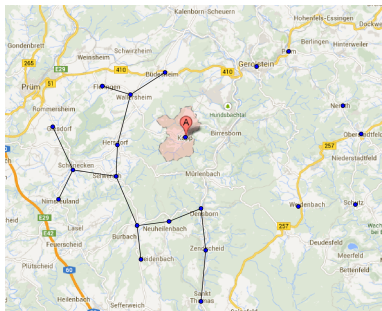


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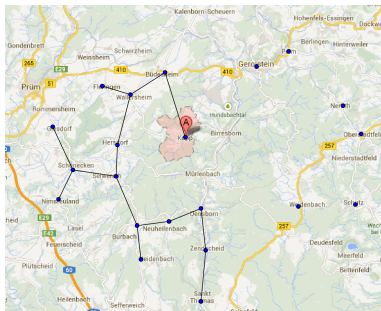


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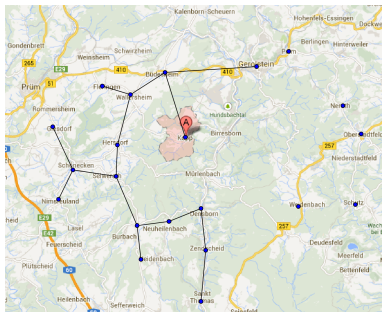
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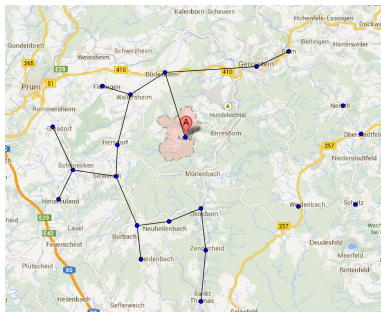


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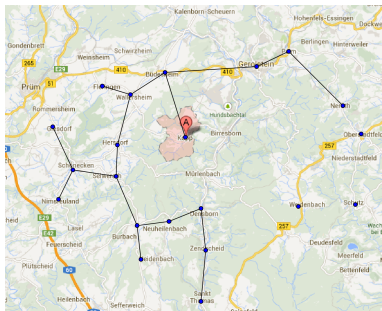


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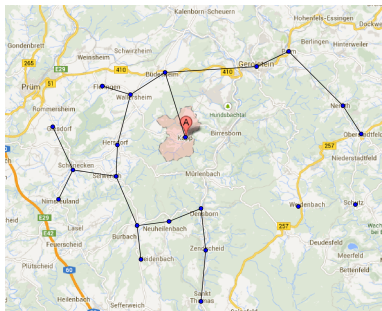


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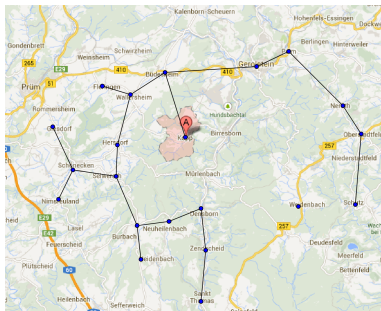


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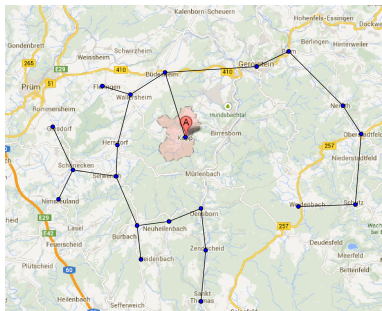


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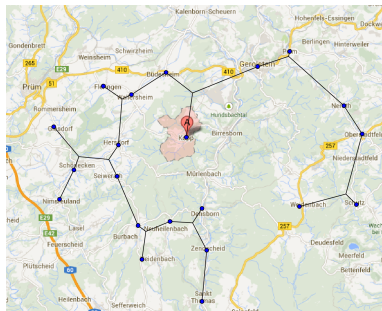
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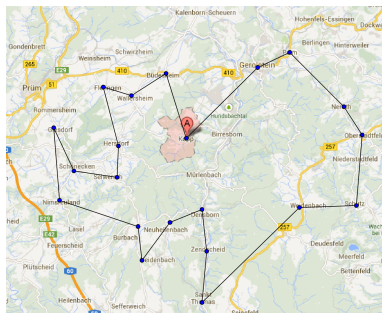
The resulting spanning tree is minimal in length. Call it  $G$ .

# The Traveling Salesman Problem

Suppose  $K$  connects all these points with minimal length.  
You can do a tour with double length...



# The Traveling Salesman Problem



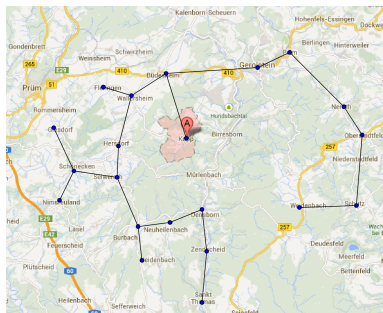
$$2\ell(K) \geq \ell(T).$$

Suppose  $K$  connects all these points with minimal length. You can do a tour with double length...

But you can shorten it by taking straight lines instead of repeating vertices. Call the minimal tour  $T$ .



# The Traveling Salesman Problem



$$2\ell(K) \geq \ell(T) > \ell(G).$$

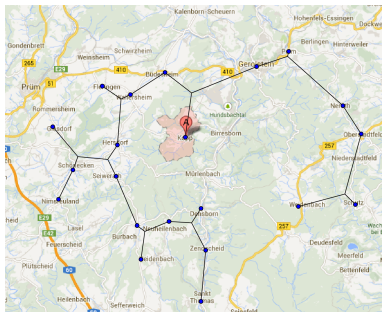
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Obviously,  $T$  contains a spanning tree and  $G$  is minimal among them.

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$$2l(K) \geq l(T) > l(G) \geq l(K):$$

Suppose  $K$  connects all these points with minimal length. You can do a tour with double length...

But you can shorten it by taking straight lines instead of repeating vertices. Call the minimal tour  $T$ .

Obviously,  $T$  contains a spanning tree and  $G$  is minimal among them.

Furthermore,  $K$  is shorter than any spanning tree.

The greedy algorithm provides us with the best route up to constant 2.

## Non-finite sets

- ▶ When it comes to a non-finite set  $E$ , the Traveling Salesman Problem consists in finding a minimal rectifiable curve  $\Gamma \supset E$ . This would give us also a minimal tour up to a constant.

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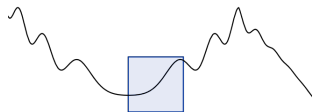
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- ▶ We say that a set is rectifiable when the set is contained in the image of a finite interval by a Lipschitz function.
- ▶ One necessary condition for  $E$  to be rectifiable is that the Hausdorff one-dimensional (outer) measure of the set,  $\mathcal{H}^1(E)$ , is finite, but it is not sufficient unless  $E$  is connected [Falconer].

# The Peter Jones' Betas (1)

## Definition

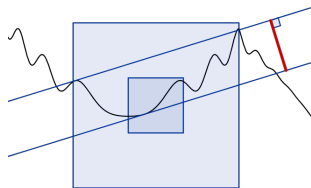
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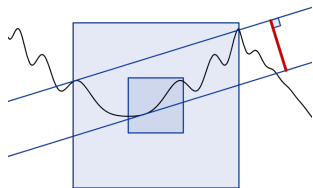
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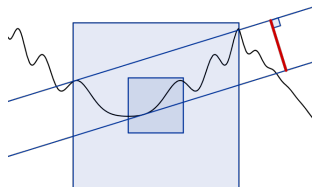
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Notice that  $\beta_E(Q) \leq 1$ . We also have that

$$\beta_E(Q) = \frac{2}{\ell(3Q)} \inf_L \{ \sup_{E \cap 3Q} \text{dist}(z, L) \}.$$



# The Peter Jones' Betas (2)

## Definition

Given a set  $E$ , we associate to it the coefficient

$$\beta^2(E) = \text{diam}(E) + \sum_{\substack{Q \in \Delta \\ \ell(Q) \leq \text{diam}(E)}} \beta_E^2(Q) \ell(Q).$$

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Notice that  $\beta_E(Q)$  is adimensional, and we can see that  $\beta^2$  has a linear behavior up to a constant.

# Main Result

## Theorem

Suppose  $E \subset \mathbb{C}$  is a bounded set. Then  $E$  is contained in a rectifiable curve if and only if  $\beta^2(E)$  is finite. Moreover, there are constants  $c_1, c_2$  such that

$$c_1 \beta^2(E) \leq \inf_{\Gamma \supset E} \mathcal{H}^1(\Gamma) \leq c_2 \beta^2(E) \quad (1)$$

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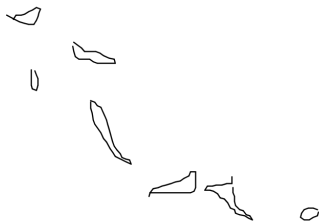
Notice that, even though we do not find the best path for the salesman, we bound the distance the salesman must travel if he designs his route wisely.

# Finding a good route

We are about to prove that  $\beta^2(E) < \infty$  implies that  $E$  is rectifiable, with

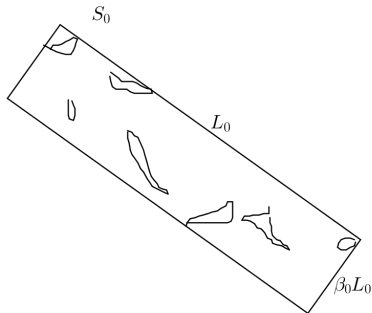
$$\inf_{\Gamma \supset E} \mathcal{H}^1(\Gamma) \leq c_2 \beta^2(E).$$

## Breaking a rectangle into two



- ▶ Consider a given bounded set  $E_0$ .

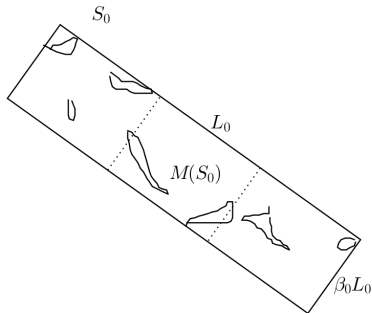
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- ▶ Consider a given bounded set  $E_0$ .
- ▶ Cover it by a strip of minimal width and shrink to a rectangle  $S_0$ .

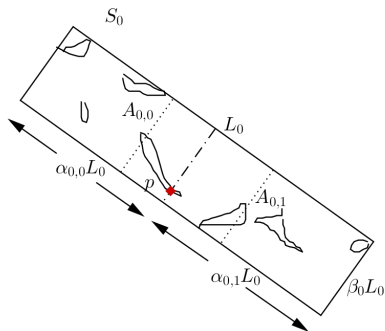


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- ▶ Consider a given bounded set  $E_0$ .
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- ▶ Divide in three equal parts.

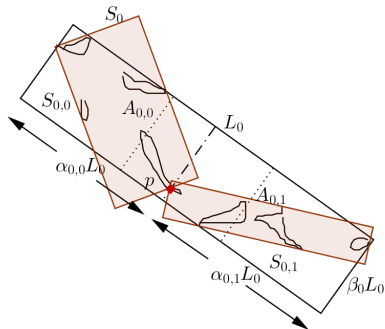
# Breaking a rectangle into two



Case 1: The middle third has a point  $p \in E_0$ .

- ▶ Divide  $S_0$  in two rectangles  $A_{0,j}$  by  $p$ .
- ▶  $\alpha_{0,0} + \alpha_{0,1} = 1$ .

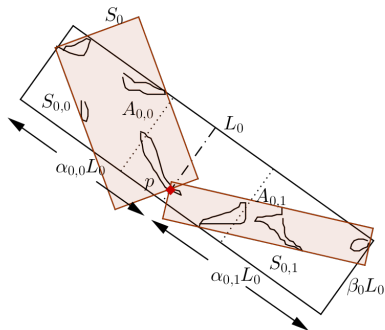
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- ▶ Cover each part  $E_{0,j}$  by a rectangle as before.

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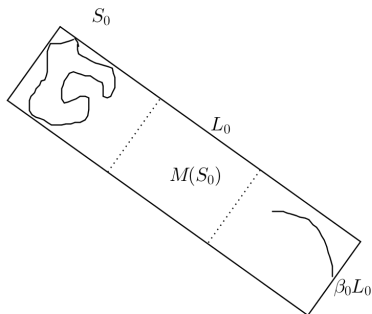


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$$L_{0,j} \leq \sqrt{\alpha_{0,j}^2 + \beta_0^2} L_0 \leq (1 + 5\beta_0^2) \alpha_{0,j} L_0.$$

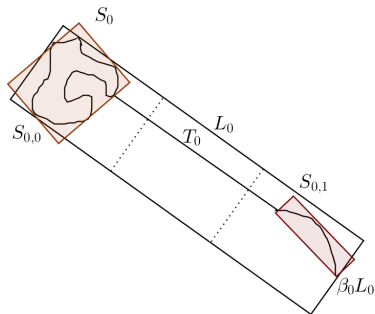
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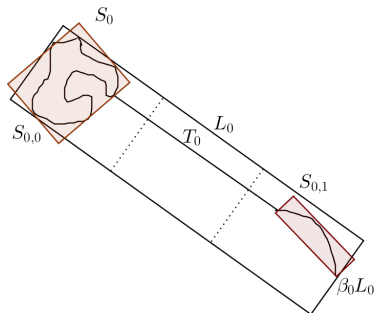
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- ▶ Still  $L_{0,j} \leq (1 + 5\beta_0^2)\alpha_{0,j}L_0$ .
- ▶ The minimal segment joining them  $|T_0| \leq (1 + \beta_0^2)L_0$ .

# Iteration

- ▶ Cover  $E$  by a rectangle  $S_0$  as before. Break the rectangle as shown to get  $S_{0,0}$  and  $S_{0,1}$ . In case 2 you get also a segment  $T_0$ .



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- ▶ Iterate the process as usual. After  $n$  steps, you have  $2^n$  rectangles  $S_I$  covering  $E_I$ ,  $I = (0, i_1, \dots, i_n)$ , with long sides  $L_I$ , weights  $\alpha_I \in [1/3, 2/3]$  and factors  $\beta_I$ .

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- ▶ After  $N = 25$  steps the diameter of a rectangle drops by at least  $1/2$ .

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Let  $R_n$  be the sum of the diameters of the rectangles at stage  $n$ .

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## Bounds for the length (1): the rectangles

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Arguing by induction, we will have the uniform bound

$$R_n \leq R_0 + \sum_{|I| \leq n} 5\beta_I^2 L_I \lesssim \text{diam}(E) + \sum \beta_I^2 L_I.$$

## Bounds for the length (2): the segments in wide rectangles

If case 2 is applied and  $L_{I,0} + L_{I,1} \geq 0.9L_I$ , then  $\beta_I \geq \tilde{\beta}$  for a fixed constant  $\tilde{\beta}$ .

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Thus, the sum of the lengths of the middle segments created from applications of case 2 with  $L_{I,0} + L_{I,1} \geq 0.9L_I$  is at most  $\sum \beta_I^2L_I$ .

## Bounds for the length (3): the segments in narrow rectangles

Now write

$$R_n = I_n + II_n$$

where  $II_n$  is the sum of the lengths of the rectangles at stage  $n$  to which case 2 will be applied and for which  $L_{I,0} + L_{I,1} < 0.9L_I$ . Let  $T_{n+1}$  denote the sum of the lengths of the segments created when  $L_{I,0} + L_{I,1} < 0.9L_I$  at stage  $n$ .

## Bounds for the length (3): the segments in narrow rectangles

Now write

$$R_n = I_n + II_n$$

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$$R_{n+1} \leq \left( I_n + C \sum_{|I|=n+1} \beta_I^2 L_I \right) + 0.9 II_n.$$

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Moreover, as  $|T_I| \leq (1 + \beta_I^2)L_I$ , we have

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Summing both inequalities,

$$R_{n+1} + 0.1 T_{n+1} \leq R_n + C \sum_{|I|=n+1} \beta_I^2 L_I.$$

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Equivalently,

$$0.1 T_{n+1} \leq R_n - R_{n+1} + C \sum_{|I|=n+1} \beta_I^2 L_I.$$

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As  $R_n$  is uniformly bounded by  $C \text{diam}(E) + C \sum \beta_I^2 L_I$ , also

$$\sum_n T_n \lesssim \text{diam}(E) + \sum \beta_I^2 L_I.$$

## Summing cubes

It only remains to bound

$$\sum \beta_I^2 L_I \leq \sum_Q \beta_E^2(Q) \ell(Q).$$



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Given  $I$  chose a dyadic cube  $Q_I$  such that  $d_I := \text{diam}(S_I) \leq \ell(Q_I) < 2d_I$  and with  $Q_I \cap S_I \neq \emptyset$ . As  $3Q_I \supset S_I$ , we will have the bound

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It will suffice to prove that

$$\sum_{I:Q_I=Q} \beta_I^2 L_I \leq C \beta_E^2(Q) \ell(Q).$$

## Summing cubes

Indeed, write  $\mathcal{F}(I)$  for the father index of  $I$ ,  $(0, i_1, \dots, i_{n-1})$ , and call  $J \geq I$  if  $I = \mathcal{F}^j(J)$  for some  $j \geq 0$ .

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$$|J| - |I| \leq 24.$$

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We can classify the independent branches where  $Q$  occurs as follows:

$$\{I : Q_I = Q\} = \bigcup_{I: d_{\mathcal{F}(I)} > \ell(Q) \geq d_I} \bigcup_{j=0}^{24} \{J \geq I : |J| = |I| + j \text{ and } Q_J = Q\}$$

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Then, using the previous bound, given a dyadic cube  $Q$  we have

$$\sum_{I: Q_I = Q} \beta_I^2 L_I \lesssim \beta_E(Q) \ell(Q) \sum_{j=0}^{24} \sum_{\substack{J: Q_J = Q \\ d_{\mathcal{F}^{j+1}(I)} > \ell(Q) \geq d_{\mathcal{F}^j(I)}}} \beta_J$$

## Summing cubes

Now, for  $j \leq 24$ , the domains  $E_J$  appearing in the sum

$$\sum_{\substack{J: Q_J = Q \\ d_{\mathcal{F}^{j+1}(I)} > \ell(Q) \geq d_{\mathcal{F}^j(I)}}} \beta_J$$

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are contained in disjoint convex polygons  $\tilde{S}_J$  of width  $\beta_J L_J$  and diameter comparable to  $L_J$ . One can see that

$$\beta_J L_J^2 \approx \text{Area}(\tilde{S}_J).$$



## Summing cubes

At the same time, all of them are in a strip of width  $\beta_E(Q)\ell(Q)$  and contained in  $3Q$ . The areas of the polygons are bounded in consequence by the area of this strip intersected with  $3Q$ . Thus,

$$\sum_{\substack{J:Q_J=Q \\ d_{\mathcal{F}^j(I)} > \ell(Q) \geq d_{\mathcal{F}^{j-1}(I)}}} \beta_J L_J^2 \leq C \beta_E(Q) \ell(Q)^2$$

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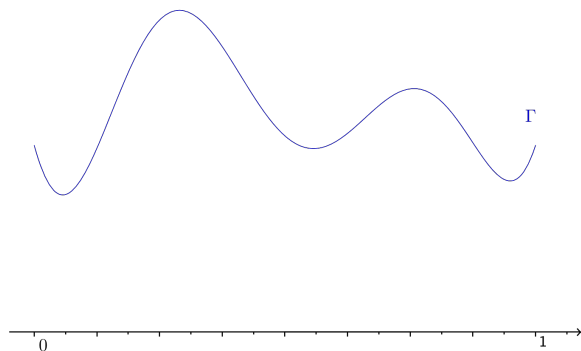
Taking into account that  $L_J \approx \ell(Q)$  and summing in  $j$  we obtain

$$\sum_{I:Q_I=Q} \beta_I^2 L_I \leq C \beta_E^2(Q) \ell(Q).$$

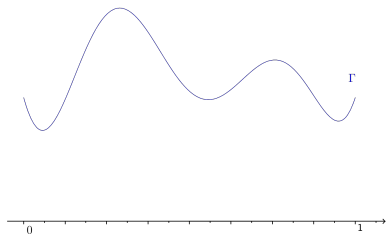
# Lipschitz graphs

## Lemma

Let  $\Gamma$  be the graph of a Lipschitz function. For  $E \subset \Gamma$ ,  $\beta^2(E) \leq C\mathcal{H}^1(\Gamma)$ .

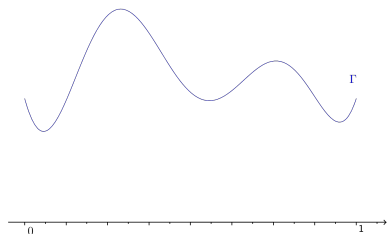


# Lipschitz graphs

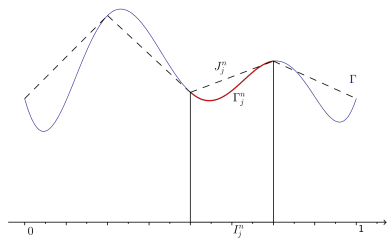


# Lipschitz graphs

Let  $\Gamma = \{0 \leq x \leq 1, y = f(x)\}$ ,  
where  $f$  is Lipschitz with constant  $M$ . It is enough to show the case  
 $E = \Gamma$  and  $f(0) = f(1)$ .



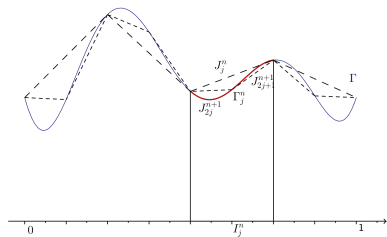
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Let  $I_j^n$  be the  $j$ th dyadic interval of length  $2^{-n}$ , call its image graph  $\Gamma_j^n$  and let  $J_j^n$  be the segment uniting the endpoints of  $\Gamma_j^n$ .

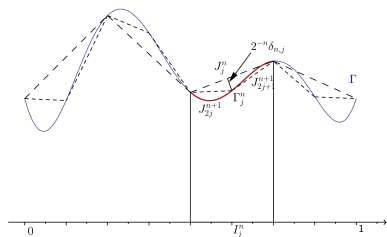
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$$2^{-n} \delta_{n,j}^2 \lesssim \ell(J_{2j}^{n+1}) + \ell(J_{2j+1}^{n+1}) - \ell(J_j^n)$$

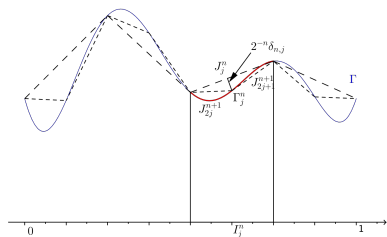
with constant depending on  $M$ .



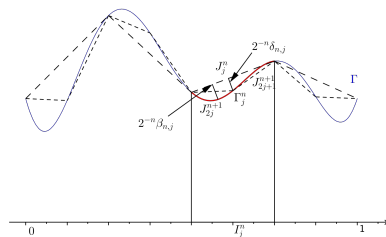
# Lipschitz graphs

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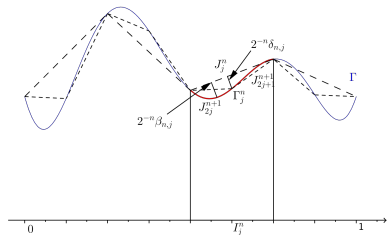
$$\sum_{m,k} c2^{-m}\delta_{m,k}^2 \leq 2\ell(\Gamma).$$

Now, by the triangular inequality,

$$\beta_{n,j} := 2^n \sup\{\text{dist}(z, J_j^n) : z \in \Gamma_j^n\},$$

$$\beta_{n,j} \leq \sum_{m=n}^{\infty} 2^{n-m} \sup\{\delta_{m,k} : I_k^m \subset I_j^n\}.$$

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Using Hölder inequalities and other standard arguments for series, one gets

$$\sum_{n,j} 2^{-n} \beta_{n,j}^2 \lesssim \sum_{m,k} 2^{-m} \delta_{m,k}^2 \lesssim 2\ell(\Gamma).$$

# Hölderizing

Indeed,

$$\sum_{n,j} 2^{-n} \beta_{n,j}^2 \leq \sum_{n,j} 2^{-n} \left( \sum_{m=n}^{\infty} 2^{n-m} \sup_{I_k^m \subset I_j^n} \delta_{m,k} \right)^2$$

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 &\leq C \sum_{n,j} \sum_{\substack{m \geq n \\ I_k^m \subset I_j^n}} 2^{\frac{n}{2}} 2^{-\frac{3m}{2}} \delta_{m,k}^2
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 &\leq C \sum_{n,j} \sum_{\substack{m \geq n \\ I_k^m \subset I_j^n}} 2^{\frac{n}{2}} 2^{-\frac{3m}{2}} \delta_{m,k}^2 \\
 &\leq C \sum_{m,k} \left( \sum_{n=0}^m \sum_{I_j^n \supset I_k^m} 2^{\frac{n}{2}} \right) 2^{-\frac{3m}{2}} \delta_{m,k}^2
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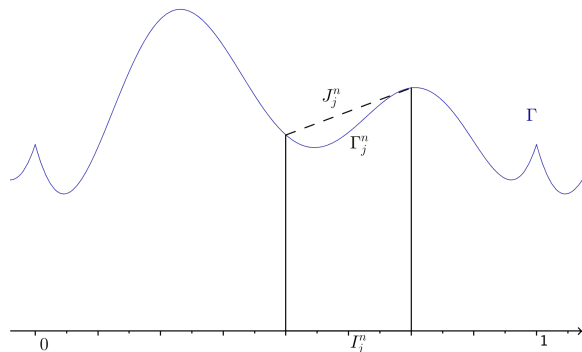
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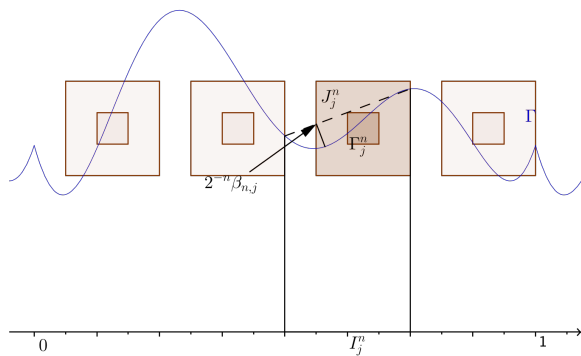
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 &\leq C \sum_{m,k} 2^{\frac{m}{2}} 2^{-\frac{3m}{2}} \delta_{m,k}^2 \lesssim \ell(\Gamma)
 \end{aligned}$$

## Final step

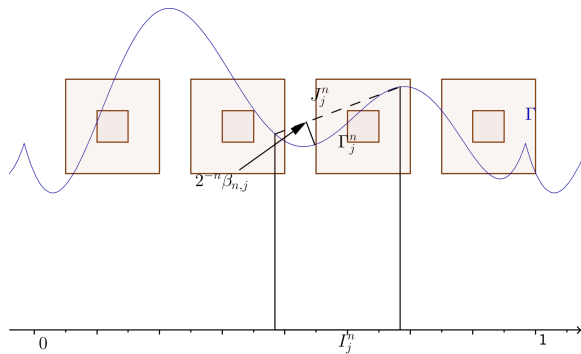
Finally we extend the function periodically and obtain some translated coefficients  $\beta_{n,j}(t)$  related to  $\Gamma(t) = (Id \times f)([t, 1+t]) \subset \mathbb{C}$  verifying the last inequality as well.



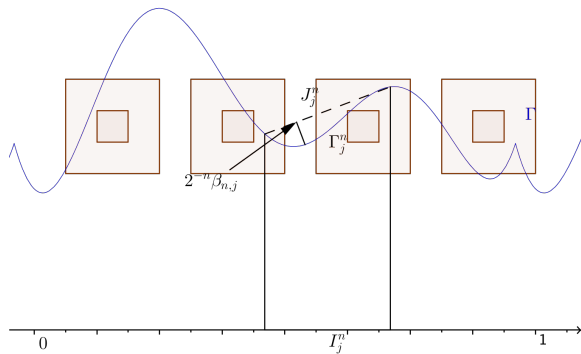
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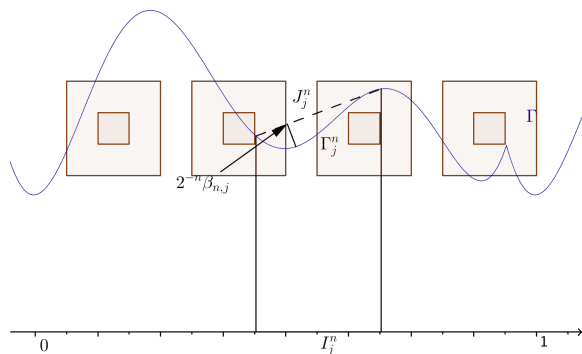
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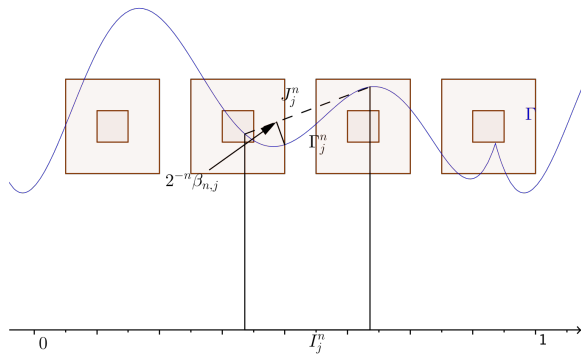
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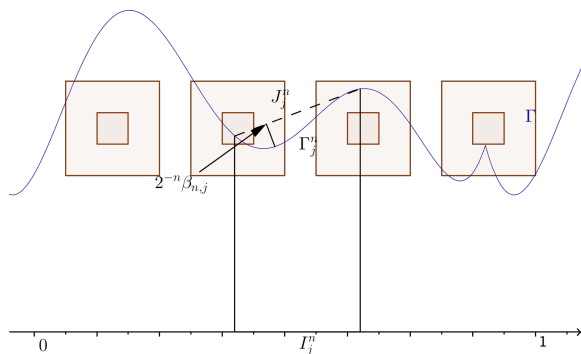
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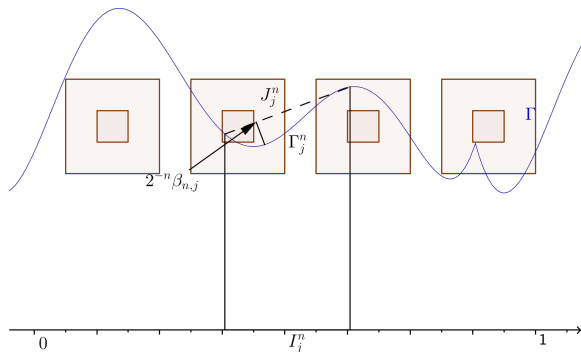


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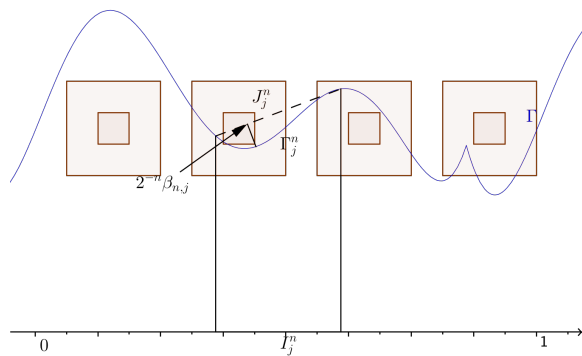




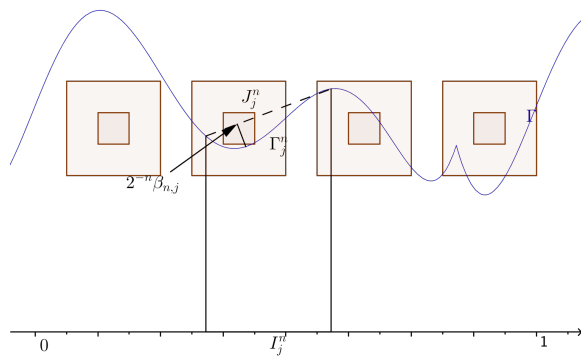
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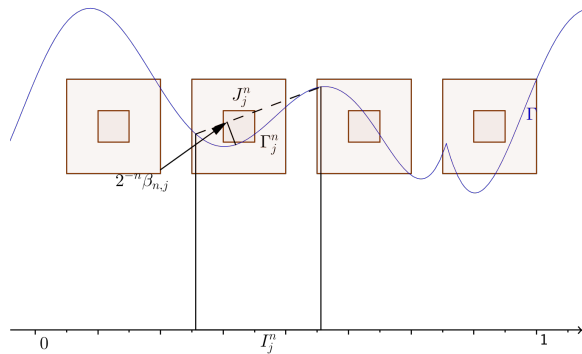
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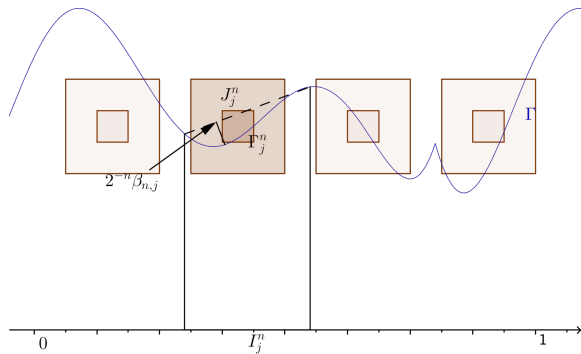
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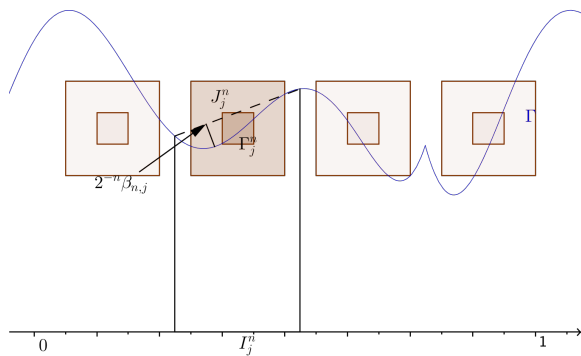
# Final step



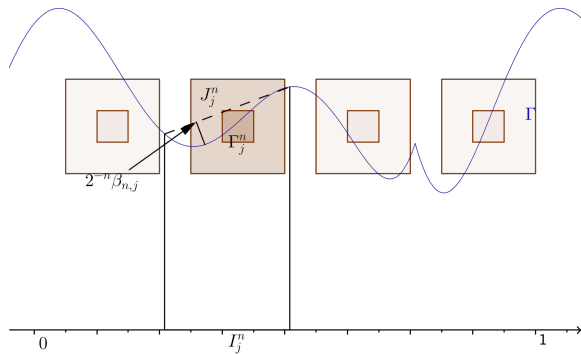
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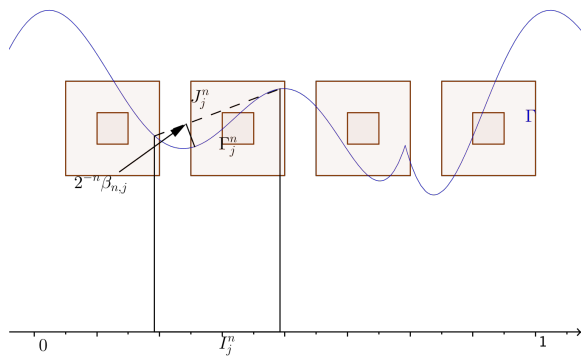
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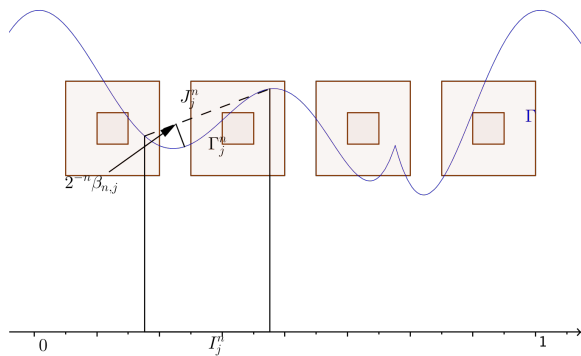


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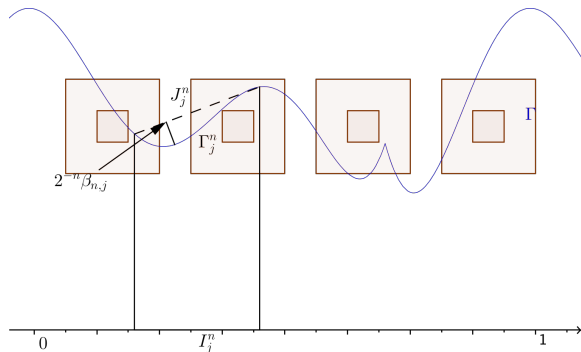




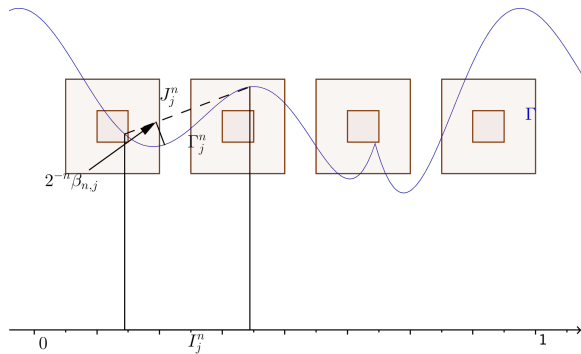
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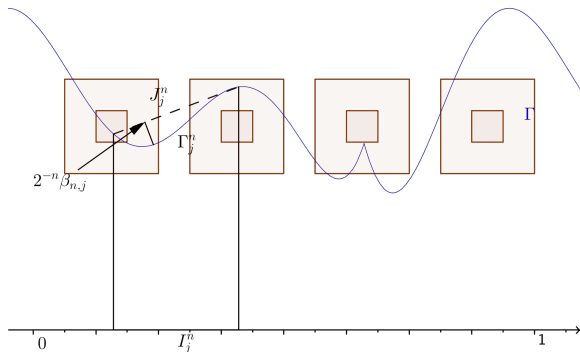
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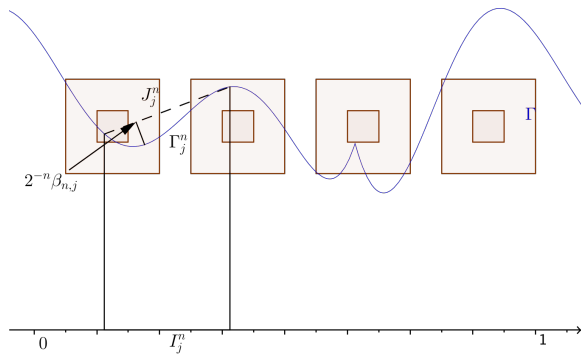
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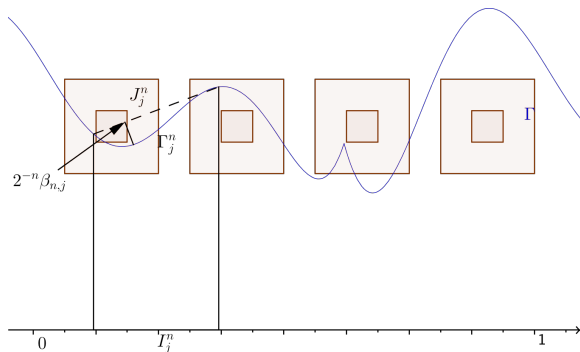
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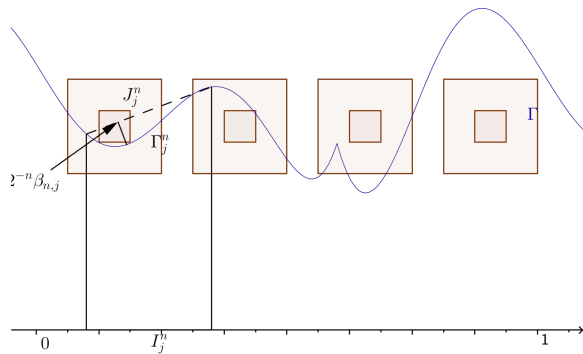
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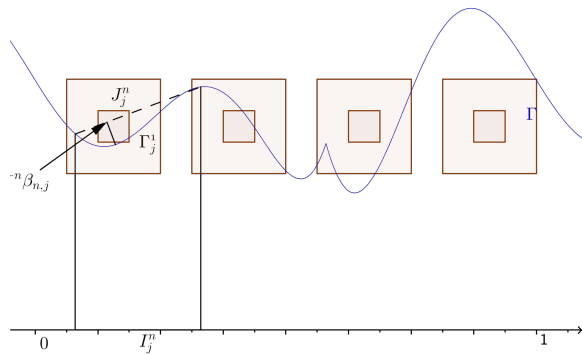
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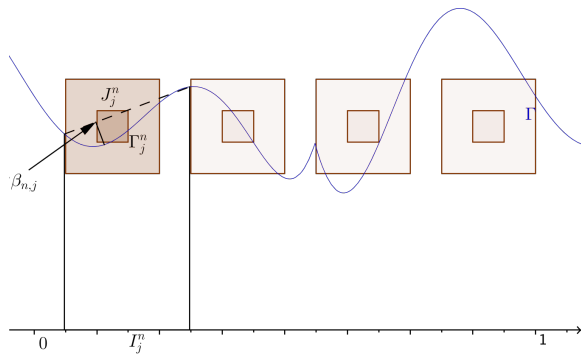


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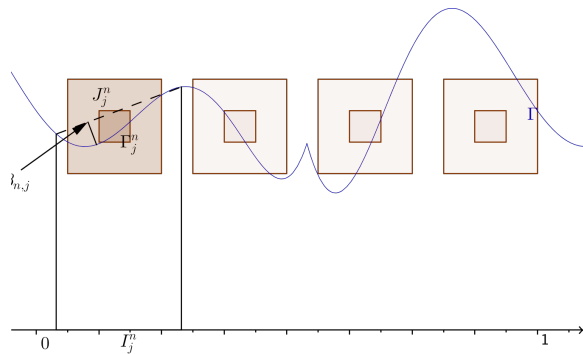




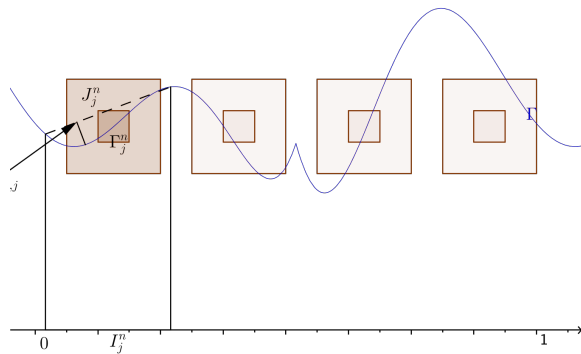
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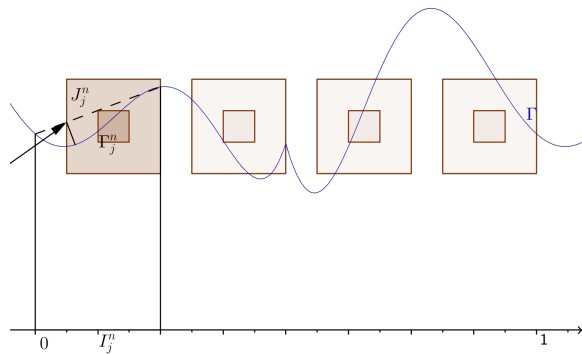
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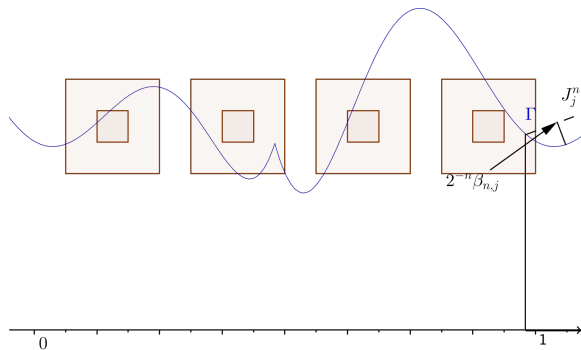
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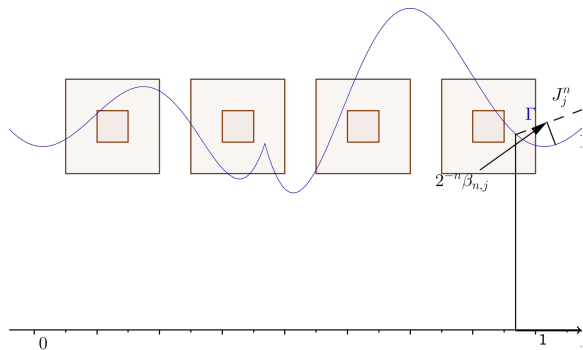
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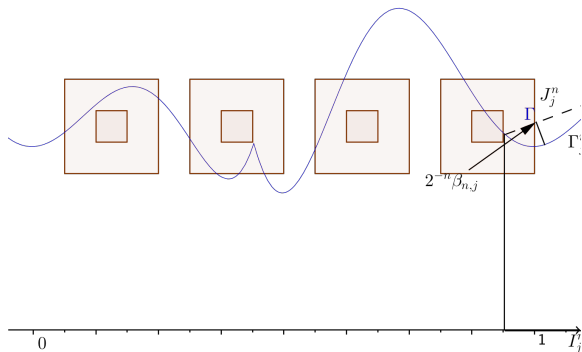
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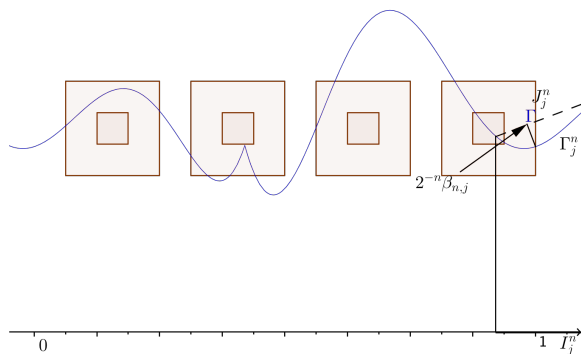
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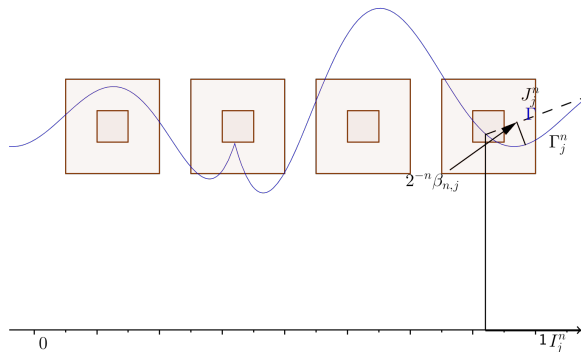


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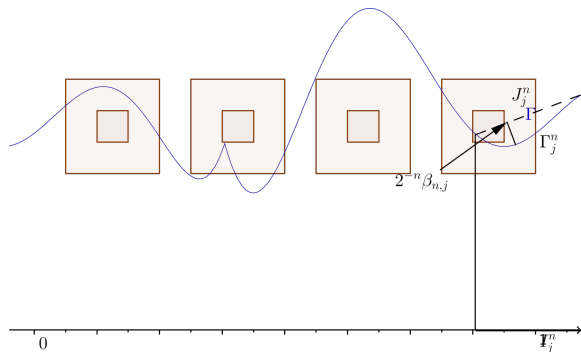




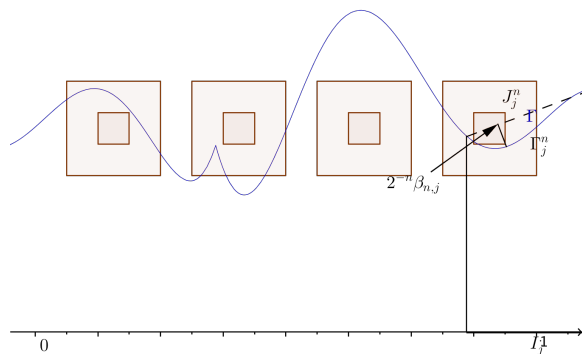
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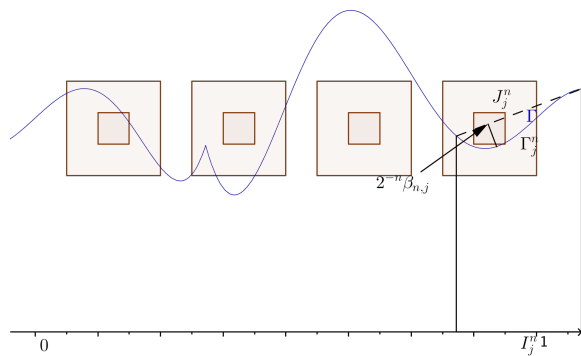
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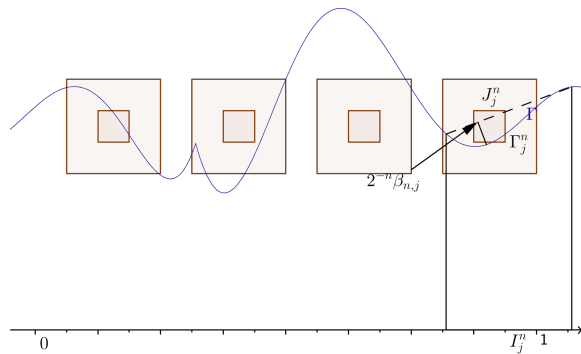
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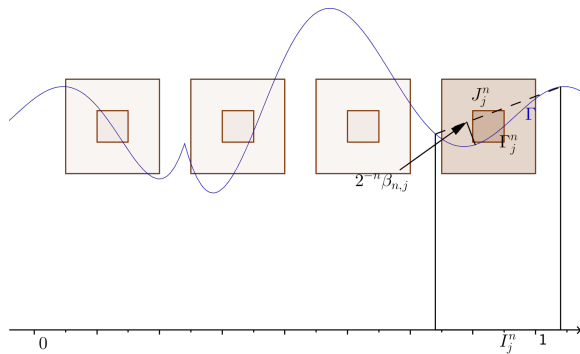
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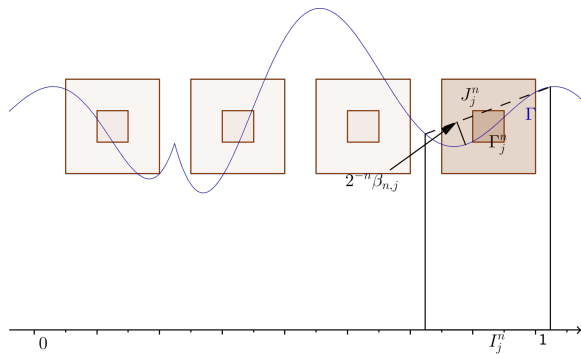
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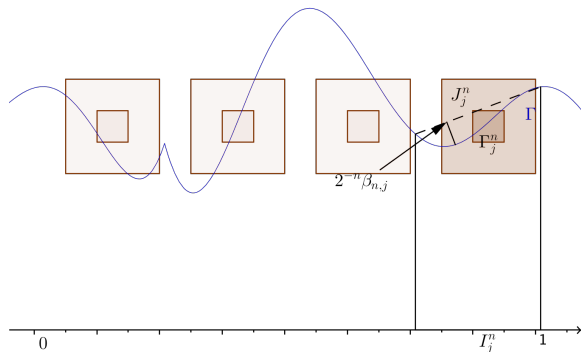
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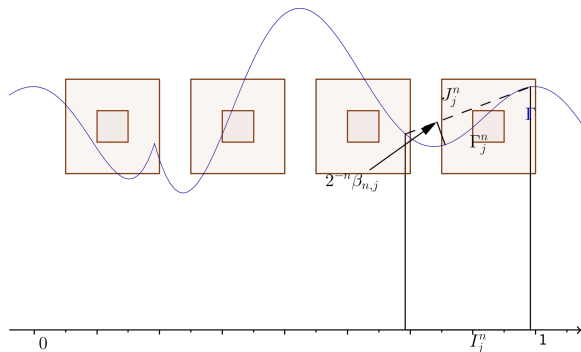


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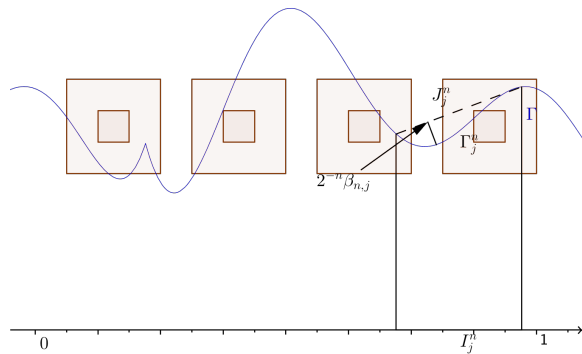




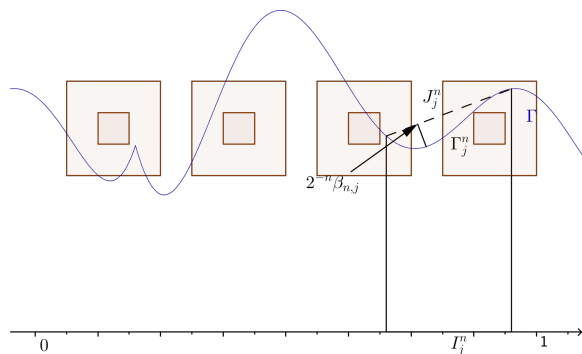
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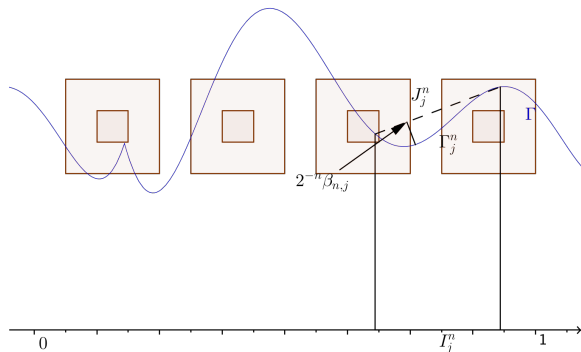
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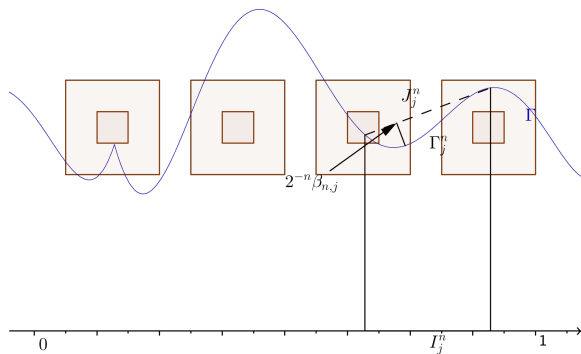
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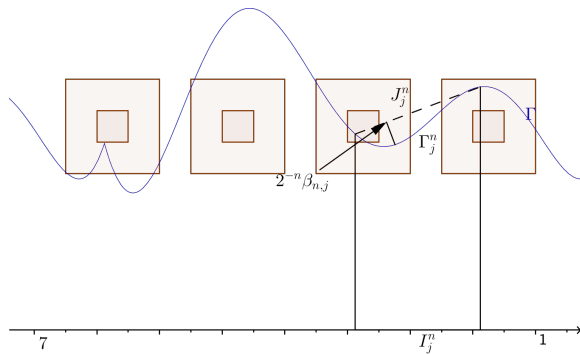
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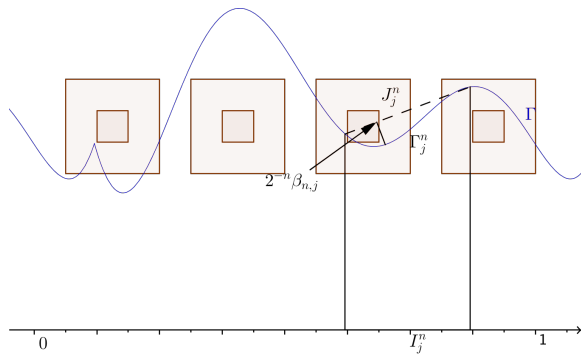
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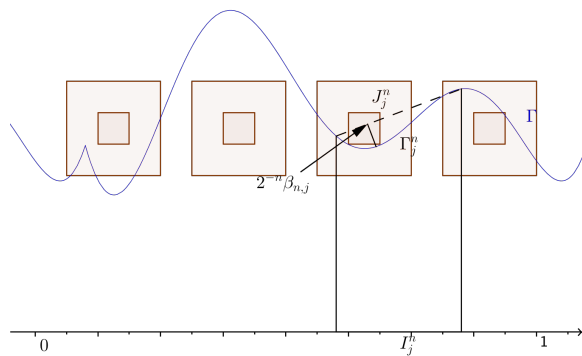
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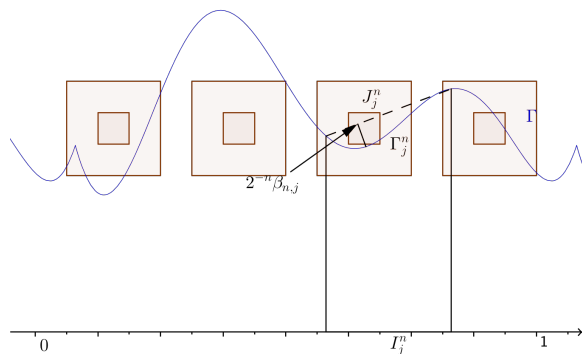


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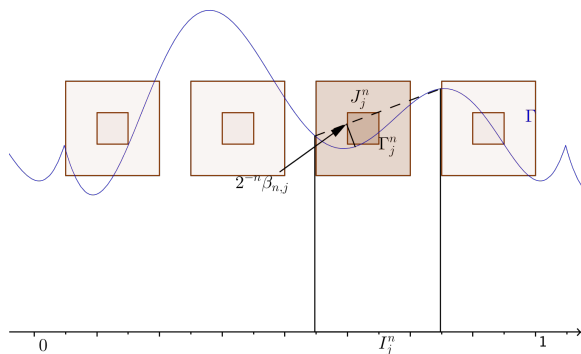




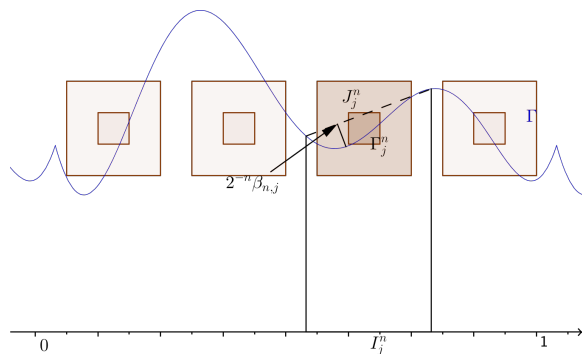
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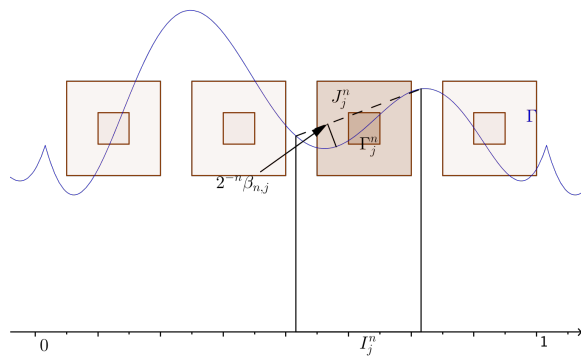
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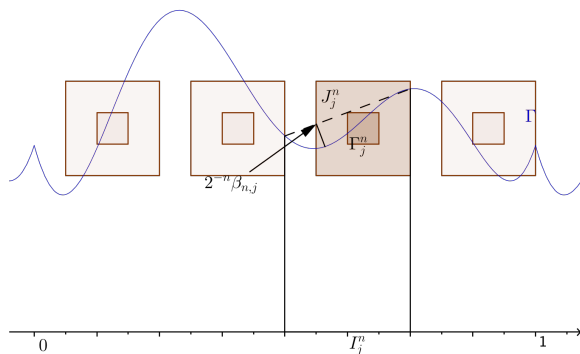
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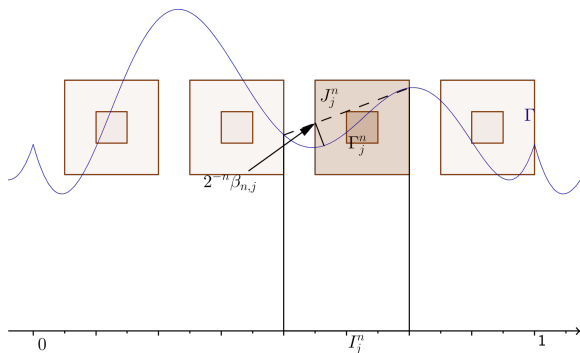
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Then, for a cube  $Q$  with  $\ell(Q) = 2^{-n-2}$ ,  $3Q$  will have projection contained in the translation of an interval  $I_{n,j}(t)$  with probability  $1/4$  with respect to the Lebesgue measure on  $t$ .



So

$$\sum_{\ell(Q)=2^{-n-2}} \beta_{\Gamma}^2(Q) \lesssim \int_{-1}^1 \sum_j \beta_{n,j}(t)^2 dt.$$

Summing with respect to  $n$  proves the claim for Lipschitz graphs.

## Second session

## Back to previous steps: Main Result

### Theorem

Suppose  $E \subset \mathbb{C}$  is a bounded set. Then  $E$  is contained in a rectifiable curve if and only if  $\beta^2(E)$  is finite. Moreover, there are constants  $c_1, c_2$  such that

$$c_1 \beta^2(E) \leq \inf_{\Gamma \supset E} \mathcal{H}^1(\Gamma) \leq c_2 \beta^2(E) \quad (2)$$

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where the infimum is taken over all rectifiable curves containing  $E$ .

We have already proven left-hand side for Lipschitz graphs and right-hand side for general sets with finite  $\beta$ .

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The key point is to find  $E \subset \bigcup \Gamma_j$  being each  $\Gamma_j$  the boundary of a Lipschitz domain  $\mathcal{D}_j$  with some restrictions on the constant and the shapes. We need to do this in such a way that we keep control on the total length and the relations between the original betas and  $\sum_Q \sum_{\Gamma_j} \beta_{\Gamma_j}^2(Q)$ .

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$$\sum_Q \sum_{\Gamma_j} \beta_{\Gamma_j}^2(Q).$$

We will do that in three steps. First we present a theorem which will allow us to make the decomposition as long as  $E$  is the boundary of a simply connected domain. The second step is a simple corollary allowing us to make such a decomposition on any connected plain set  $\gamma$ . Finally we will prove the relation between betas.

# $M$ -Lipschitz domains

## Definition

We call an  $M$ -Lipschitz domain to a simply connected domain whose boundary can be expressed as  $\{r(\theta)e^{i\theta} : 0 \leq \theta < 2\pi\}$  (i.e. it is starlike with respect to the origin), with  $r$  a Lipschitz function of coefficient  $M$  and  $\frac{1}{M+1} \leq r(\theta) \leq 1$  after translation and dilation if necessary.

# Decomposition theorem

## Theorem

*There is a constant  $M$  such that whenever  $\Omega$  is a simply connected domain with  $\mathcal{H}^1(\partial\Omega) < \infty$  there exists a rectifiable curve  $\Gamma$  such that*

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- ▶ each  $\Omega_j$  is an  $M$ -Lipschitz domain,
- ▶ and  $\sum_j \mathcal{H}^1(\partial\Omega_j) \leq M\mathcal{H}^1(\partial\Omega)$ .



## Summary of the proof

Let  $\varphi : \mathbb{D} \rightarrow \Omega$  be a Riemann mapping. By translating, rotating and rescaling the domain, we can assume WLOG that  $\varphi(0) = 0$  and  $\varphi'(0) = 1$  (i.e.  $\varphi \in S$ ).

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On the first step we will create uniformly chord-arc domains such that we keep control on the lengths. After that we will decompose these domains into smaller domains to ensure the  $M$ -Lipschitz condition is satisfied.

## Useful theorems

### Theorem (Koebe's estimate, growth and distortion theorem)

Given a conformal mapping  $\varphi \in S$  ( $\varphi : \mathbb{D} \rightarrow \Omega$ ,  $\varphi(0) = 0$ ,  $\varphi'(0) = 1$ ), we have

- ▶  $\text{dist}(\varphi(z), \partial\Omega) \approx |\varphi'(z)|(1 - |z|^2)$ .
- ▶ *Whitney cubes are almost invariant, with constant derivative absolute value on them.*
- ▶  $\frac{|z|}{(1+|z|)^2} \leq |\varphi(z)| \leq \frac{|z|}{(1-|z|)^2}$ .
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### Theorem (F. and M. Riesz Theorem)

Given a Riemann mapping  $\varphi$  to a Jordan domain  $\Omega$ , it is bounded by a rectifiable curve if and only if  $\varphi' \in H^1$ , with

$$\mathcal{H}^1(\Gamma) = \|\varphi'\|_{H^1}.$$

## Useful theorems

### Theorem (Alexander's)

For  $\Gamma$  connected with finite length, call  $\varphi_i$  to a collection of Riemann mappings to each component of  $\mathbb{C}^*$ . Then

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Call the cone  $\Gamma_\alpha(\psi) = \{z \in \mathbb{D} : |z - \psi| < \alpha(1 - |z|)\}$  and the area function  $A_\alpha\varphi(\psi) = \left(\iint_{\Gamma_\alpha(\psi)} |\varphi'(z)|^2\right)^{1/2}$ .

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### Theorem (M. Calderon's)

Let  $\Omega$  be chord-arc domain,  $\alpha > 1$ ,  $0 < p < \infty$ ,  $\varphi : \Omega$  analytic. Then

$$\|\varphi - \varphi(z_0)\|_{H^p(\Omega)}^p \approx \|A_\alpha\varphi\|_{L^p(\partial\Omega)}^p.$$



## Some tools

Write  $F = \sqrt{\varphi'}$  and  $g = \log(\varphi')$ .

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$$\|g\|_{\mathcal{B}} \leq 6 \quad (3)$$

i.e.  $\frac{|\varphi''(z)|}{|\varphi'(z)|} \leq \frac{6}{1-|z|^2}$  for all  $z \in \mathbb{D}$ .

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A simple computation shows that  $4F'(z)^2 = \varphi'(z)g'(z)^2$ .

## Some help from Hardy spaces

This implies that

$$\iint_{\mathbb{D}} |\varphi'(z)| |g'(z)|^2 \log \frac{1}{|z|} dm(z) = 4 \iint_{\mathbb{D}} |F'(z)|^2 \log \frac{1}{|z|} dm(z).$$

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By the Littlewood-Paley formula for the Hardy space  $H^2$ , we have

$$\iint_{\mathbb{D}} |\varphi'(z)| |g'(z)|^2 \log \frac{1}{|z|} dm(z) \leq 2 \|F\|_{H^2}^2 = 2 \|\varphi'\|_{H^1}^2.$$

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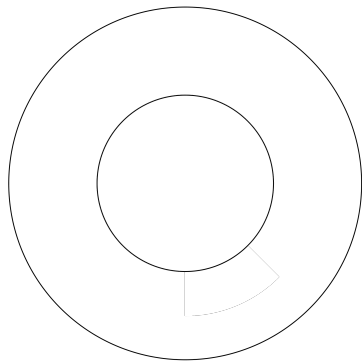
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Thanks to a result due to Alexander (which somehow generalizes the F. and M. Riesz Theorem to any simply connected domain) we can see that  $\varphi' \in H^1$ , and  $\|\varphi'\|_{H^1} \leq 2\mathcal{H}^1(\partial\Omega)$ . Summing up,

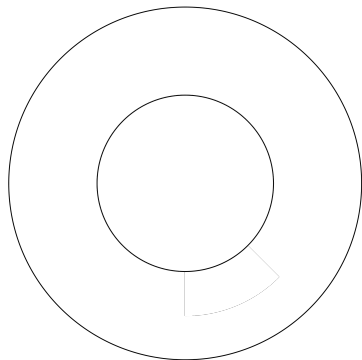
$$\iint_{\mathbb{D}} |\varphi'(z)| |g'(z)|^2 \log \frac{1}{|z|} dm(z) \leq 4\mathcal{H}^1(\partial\Omega).$$

# The local zone



Set  $\mathcal{D}_0 = \{|z| \leq 1/2\}$  and  
 $\mathcal{U}_0 = \varphi(\mathcal{D}_0)$ .

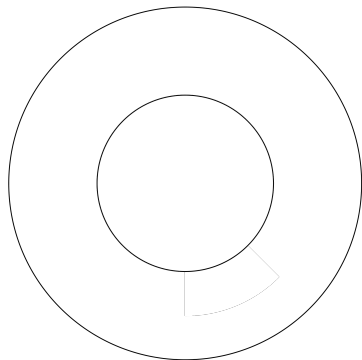
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Set  $\mathcal{D}_0 = \{|z| \leq 1/2\}$  and  $\mathcal{U}_0 = \varphi(\mathcal{D}_0)$ . By the growth theorem and the distortion theorem for univalent functions, one can see that  $\mathcal{U}_0$  is an  $M$ -Lipschitz domain.



## The local zone



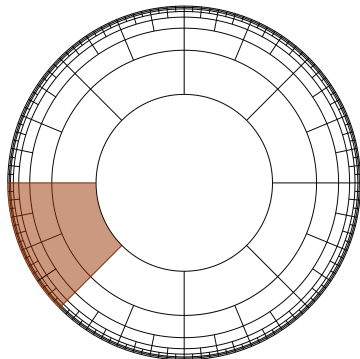
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Since  $\varphi' \in H^1$  we also have

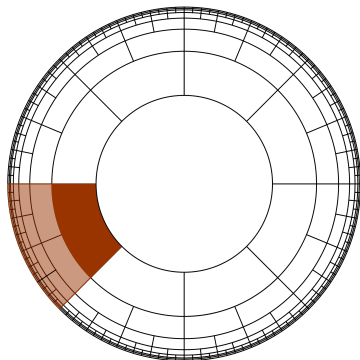
$$\mathcal{H}^1(\partial\mathcal{U}_0) \leq \mathcal{H}^1(\partial\Omega). \quad (4)$$

## Carleson boxes

Next form the dyadic Carleson boxes

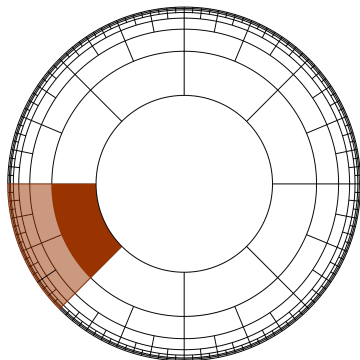


## Carleson boxes



Next form the dyadic Carleson boxes and consider their top halves  $T(Q) = \{z \in Q : |z| < 1 - 2^{-(n+1)}\}$ . Write  $z_Q$  for the center of  $T(Q)$ .

## Carleson boxes

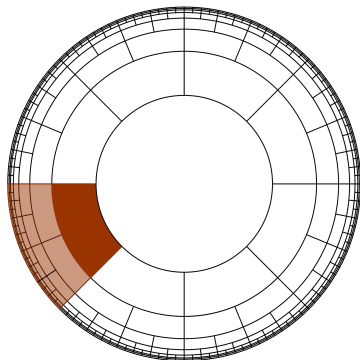


Next form the dyadic Carleson boxes and consider their top halves  $T(Q) = \{z \in Q : |z| < 1 - 2^{-(n+1)}\}$ . Write  $z_Q$  for the center of  $T(Q)$ .

We will choose the domains by a stopping time argument.

The domains  $\mathcal{D}_j$  will be unions of  $T(Q)$  so that we have a covering of the unit disk with disjoint interiors.

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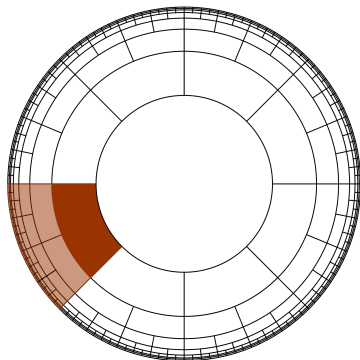


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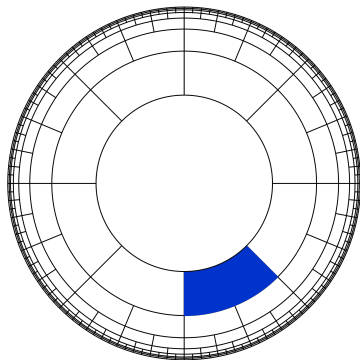
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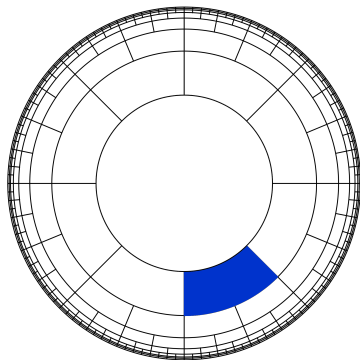
Finally we will make a subdivision of those domains to get starlike domains with uniform Lipschitz constant.

## Type 0 cubes



Fix  $\varepsilon$  to be determined later and consider a Carleson box  $Q$  as big as possible.

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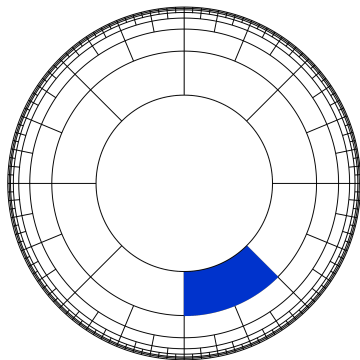
If

$$\sup_{T(Q)} |g(z) - g(z_Q)| \geq \varepsilon,$$

we say that  $Q$  is a type 0 cube and define  $\mathcal{D}(Q) = T(Q)$ .



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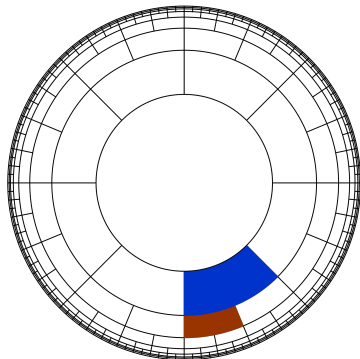
## Stopping time argument: almost constant derivative

If  $Q$  is not of type 0, define  $G(Q)$  to be the set of maximal boxes  $Q' \subset Q$  for which

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and define

$$D(Q) = \left( Q \setminus \bigcup_{G(Q)} Q' \right).$$

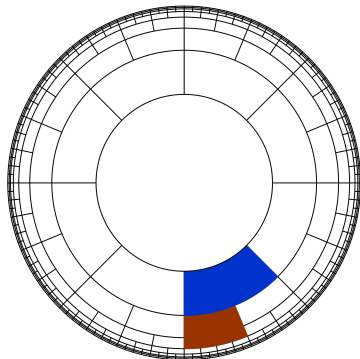


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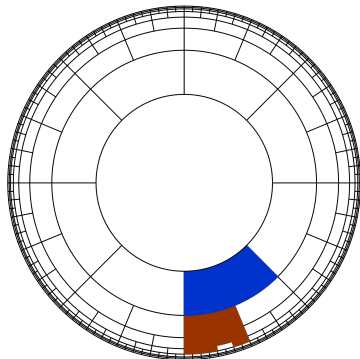
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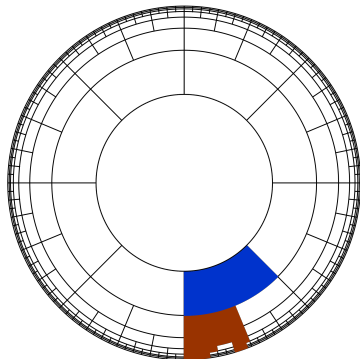
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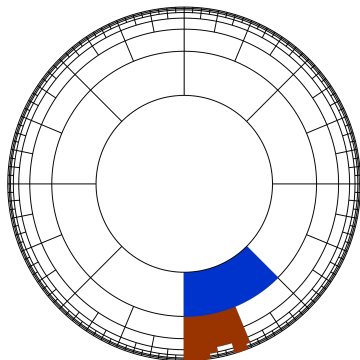
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Then,  $\mathcal{D}(Q)$  is a chord-arc domain with constant 4 and

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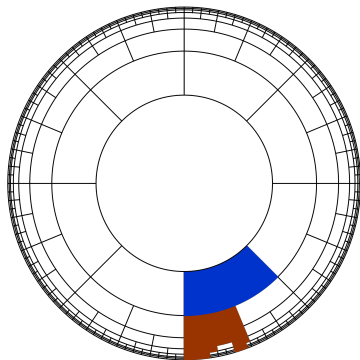
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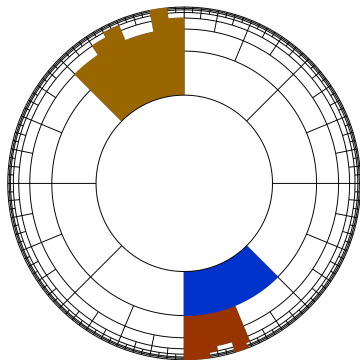
For  $\varepsilon$  small enough,  $\mathcal{U}_Q = \varphi(\mathcal{D}_Q)$  will be chord-arc domains with constant 5.

## Type 1 and type 2



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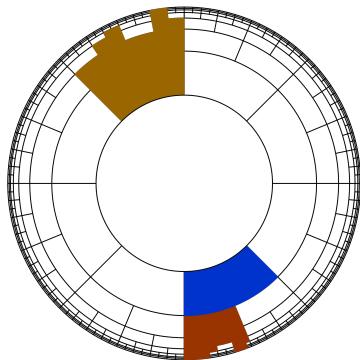


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Keep finding  $\mathcal{D}(Q)$  for the successive remaining maximal cubes in  $Q \setminus \mathcal{D}(Q)$ . Then, the family  $\{\mathcal{D}_j\}_{j \geq 0}$  is pairwise disjoint.

## Lengths in domains of type 0

If  $Q$  is of type 0, then using the Bloch norm of  $g$  and  $\sup_{T(Q')} |g(z) - g(z_Q)| \geq \varepsilon$ , we see that there is a significative part of  $T(Q)$  with  $|g'| > C\varepsilon \ell(Q)$ , so  $\ell(Q)^2 \lesssim \int |g'|^2$ .

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Finally, using the previous estimates on the last integral over  $\mathbb{D}$ ,

$$\sum_{\text{type 0}} \mathcal{H}^1(\partial \mathcal{U}_j) \leq C \mathcal{H}^1(\partial \Omega).$$

## Starlike domains come out

Dividing the region into a fixed number of polar rectangles, we can apply yet the previous reasoning. Furthermore, using again the Bloch estimate for  $g$  we find that the derivative is almost constant so that the image of the regions are  $M$ -Lipschitz domains.

## Lengths in domains of type 1

For  $Q$  of type 1, using F. and M. Riesz Theorem for Jordan domains we know that

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Finally, as this arcs have zero superposition in  $\mathcal{H}^1$ , we have using Alexander's result that

$$\sum_{\text{type 1}} \mathcal{H}^1(\partial U_j) \leq C \iint_{\partial \mathbb{D}} |\varphi'(z)| \leq C \mathcal{H}^1(\partial \Omega).$$

## Lengths in domains of type 2

When it comes to type 2 cubes, the reasoning is more involved. We sketch the proof.

Call  $\{J_k\}$  to the top edges of the boxes in  $G(Q)$ . Then

$$\mathcal{H}^1(J_k) \geq \frac{\mathcal{H}^1(\partial\mathcal{D}(Q))}{12}.$$

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This allows us to prove that

$$\mathcal{H}^1(\partial\mathcal{U}_j) \lesssim \int_{\partial\mathcal{D}(Q)} |F(z) - F(z_Q)|^2.$$

## Lengths in domains of type 2

By M. Calderon's Theorem,

$$\int_{\partial\mathcal{D}(Q)} |F(z) - F(z_Q)|^2 \leq C \iint_{\mathcal{D}(Q)} |F'(z)|^2 \mathcal{H}^1(B(z, 2\text{dist}(z, \partial\mathcal{D}(Q))))$$

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and using that chord-arc domains are bounded by Ahlfors-regular curves,

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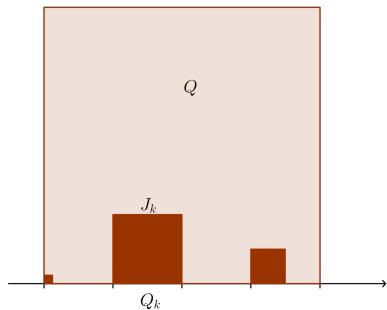
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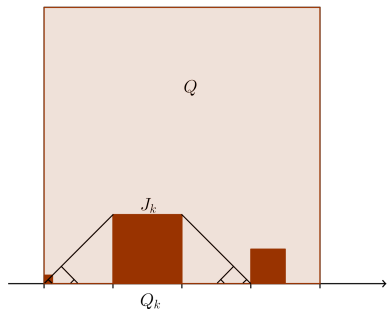
## Starlike domains in type 1 or 2

It only remains to subdivide the domains  $\mathcal{D}(Q)$  related to cubes of type 1 and type 2 into domains  $\mathcal{D}_{Q,k}$  such that  $\mathcal{U}_{Q,k} = \varphi(\mathcal{D}_{Q,k})$  are  $M$ -Lipschitz domains with length still bounded.

## Starlike domains in type 1 or 2: visual explanation

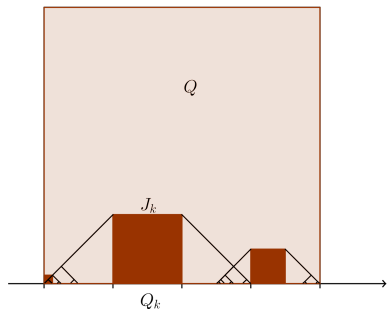


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$$\mathcal{H}^1(T(Q_k)) = Cl(Q_k)$$

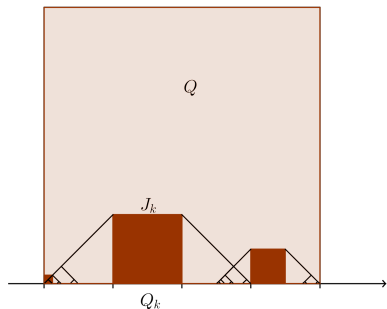
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$$\mathcal{H}^1(T(Q_k)) = C\ell(Q_k)$$

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$$\sum \mathcal{H}^1(\partial\varphi(\mathcal{D}_{j,k})) \leq C|\varphi'(z_Q)|\ell(Q) \leq C\mathcal{H}^1(\mathcal{U}_j)$$

# Given any connected set

## Corollari

*There exists a constant  $M < \infty$  such that if  $\Gamma$  is a connected plane set with  $\mathcal{H}^1(\Gamma) < \infty$ , then there exists a connected plane set  $\tilde{\Gamma} \supset \Gamma$  such that  $\mathcal{H}^1(\tilde{\Gamma}) \leq M\mathcal{H}^1(\Gamma)$ , the bounded components  $\mathcal{D}_j$  of  $\mathbb{C} \setminus \tilde{\Gamma}$  are  $M$ -Lipschitz domains with  $\Gamma \subset \bigcup \partial\mathcal{D}_j$ , and the boundary of the unbounded component  $\mathcal{D}_0$  of  $\mathbb{C} \setminus \tilde{\Gamma}$  is a circle at least  $3\sqrt{2}\mathcal{H}^1(\Gamma)$  units from  $\Gamma$ .*

# The shortest proof

## Proof.

Apply the previous result to each bounded component of the original set united to a circle big enough by a segment. □

## Small domains, big domains

Now, let  $\Gamma$  be connected with  $\mathcal{H}^1(\Gamma) < \infty$ , let  $\{\mathcal{D}_j\}$  be the Lipschitz domains given by the previous corollary and write  $\Gamma_j = \partial\mathcal{D}_j$  and  $\delta_j = \text{diam}(\mathcal{D}_j)$ .



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$$\mathcal{F}(Q) = \{\Gamma_j : \Gamma_j \cap 3Q \neq \emptyset, \delta_j \geq \ell(Q)\}$$

and

$$\mathcal{G}(Q) = \{\Gamma_j : \Gamma_j \cap 3Q \neq \emptyset, \delta_j < \ell(Q)\}.$$

## The relation between betas

### Lemma

There is a constant  $C$  such that if  $\ell(Q) \leq \text{diam}\Gamma$  and  $\ell(Q) = \frac{1}{4}\ell(Q')$ , with  $Q \subset Q'$ , then

$$\beta_{\Gamma}^2(Q) \leq C \sum_{\mathcal{F}(Q)} \beta_{\Gamma_j}^2(Q') + C_1 \frac{1}{\ell(Q)^2} \sum_{\mathcal{G}(Q)} \text{Area}(\mathcal{D}_j).$$

# Trivialities

WLOG, WMA that  $\ell(Q) = 1$  and  $\beta_{\Gamma}(Q) > 0$ , so that  $3Q \cap \Gamma_j \neq \emptyset$  for some  $\Gamma_j$  and  $3Q' \subset \bigcup \mathcal{D}_j$ .

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If  $\mathcal{F}(Q) = \emptyset$  then  $\sum_{\mathcal{G}(Q)} \text{Area} \mathcal{D}_j \geq 9\ell(Q)^2$ . Thus, we can assume there exists  $\Gamma_1 \in \mathcal{F}(Q)$ .

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We distinguish three cases.

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Hence

$$\beta_{\Gamma}^2(Q) \leq (d + d_0)^2 \leq 2d^2 + 2d_0^2 \lesssim \beta_{\Gamma_1}^2(Q') + \sum_{g(Q)} \text{Area} \mathcal{D}_j.$$

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Consequently,

$$\beta_{\Gamma}^2(Q) \lesssim \beta_{\Gamma_1}^2(Q') + \beta_{\Gamma_2}(Q')^2 + d(z_3, \Gamma_j)^2 \lesssim \sum_{\mathcal{F}(Q)} \beta_{\Gamma_j}^2(Q') + \sum_{\mathcal{G}(Q)} \text{Area}(\mathcal{D}_j)$$



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Then, because each  $\mathcal{D}_j$  is an  $M$ -Lipschitz domain, there exist at least one  $\Gamma_j \in \mathcal{F}(Q)$  such that  $\beta_{\Gamma_j}(3Q') \geq C_1$ , as three strips intersecting  $3Q$  will always intersect one another in  $3Q'$ .

# Proof of the theorem

To finish the proof of the main theorem, let  $\Gamma$  be a rectifiable curve and let  $\{\Gamma_j\}$  be as in the corollary. Using the lemma on Lipschitz graphs we can see that

$$\sum_Q \beta_{\Gamma_j}^2(Q) \ell(Q) \leq C \ell(\Gamma_j).$$

## Small cubes' areas

If  $\delta_j < 2^{-n}$  there are at most 25 dyadic cubes  $Q$  such that  $\ell(Q) = 2^{-n}$  and  $\mathcal{D}_j \in \mathcal{G}(Q)$ .

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 &\leq C \sum_j \ell(\Gamma_j) \leq C \ell(\Gamma)
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On the other hand, in the sum

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Thank you!